

Near-optimal hierarchical matrix approximation from matrix-vector products

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`chen.pw/slides`

Collaborators

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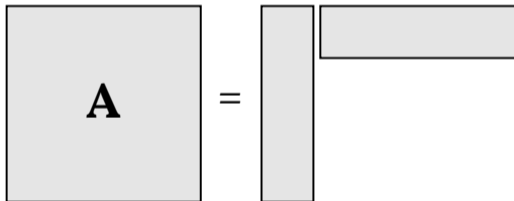


Matrix recovery and approximation

Let S be some family of matrices parameterized by a few parameters.

Matrix recovery and approximation

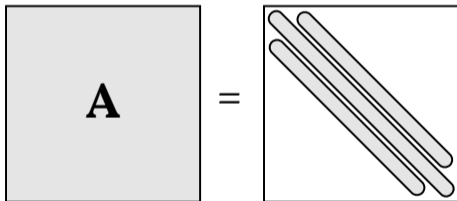
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A diagram illustrating matrix factorization. On the left is a large square matrix labeled **A**. To its right is an equals sign. Further right are two smaller matrices: a tall, narrow vertical rectangle and a wide, short horizontal rectangle, positioned such that their product equals matrix A.

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Matrix recovery and approximation

Recovery: Promised $\mathbf{A} \in S$, learn parameterization of \mathbf{A} .

- past work for many classes: low-rank, sparse, circulant, hierarchical (HODLR, HSS, etc.), butterfly, etc.¹

Approximation: Arbitrary \mathbf{A} , learn (parameterization of) $\tilde{\mathbf{A}} \in S$ such that

$$\|\mathbf{A} - \tilde{\mathbf{A}}\| \leq (1 + \varepsilon) \min_{\mathbf{X} \in S} \|\mathbf{A} - \mathbf{X}\|.$$

- lots theory on low-rank approximation, but not much else

¹Halikias and Townsend 2023.

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Access model

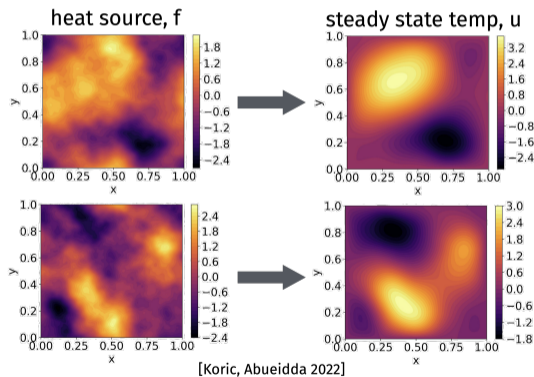
Assume we can only access \mathbf{A} using matrix-vector (matvec) queries $\mathbf{x} \mapsto \mathbf{A}\mathbf{x}$ or $\mathbf{y} \mapsto \mathbf{A}^\top \mathbf{y}$.

Why might this access model arise?

- if $\mathbf{A} = \mathbf{B}^{-1}$, we can compute $\mathbf{A}\mathbf{x}$ using a fast solver
- the action of \mathbf{A} could also correspond to some physical process

Operator Learning²

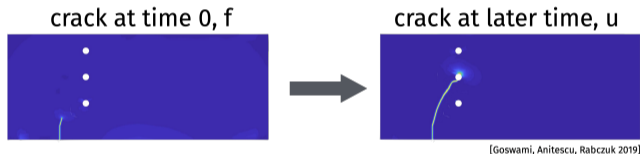
Physical processes often map a function f to a function u . I.e., implement some operator $\Phi(f) = u$.



²Boullé and Townsend 2024.

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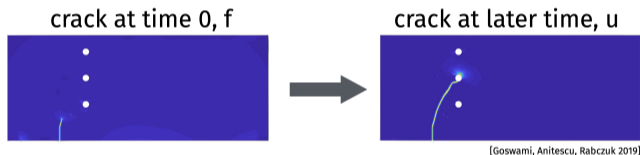
Goal: Learn mapping from input-output pairs: $(f_1, u_1), \dots, (f_m, u_m)$.

Scientific ML: Assume S is some parameterized family (e.g. neural net as in DeepONet, DeepGreen, etc.)

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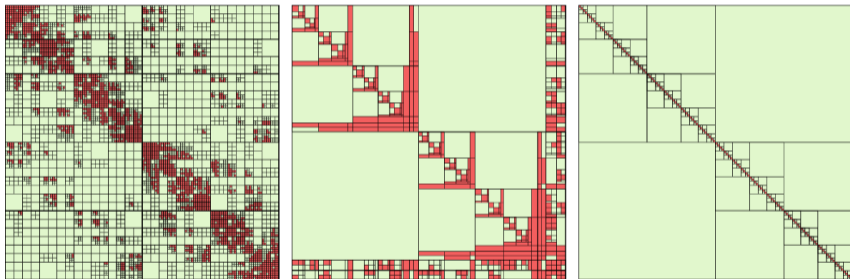
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Hierarchical matrices

Today, S will be some family of **hierarchical matrices**.



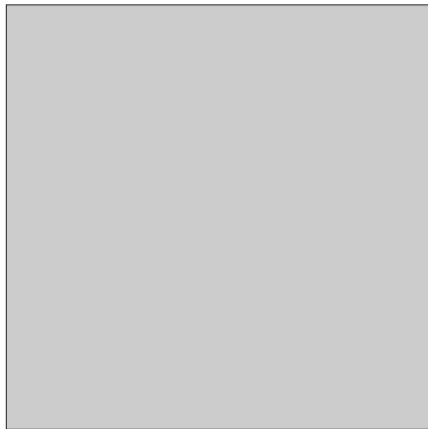
example classes: hierarchical off-diagonal low-rank (HODLR), hierarchical semi-seperable (HSS), \mathcal{H}^1 , \mathcal{H}^2 , hierarchical off-diagonal butterfly, etc.

Why hierarchical matrices?

Hierarchical matrices are useful for applications involving physical applications due to the presence of **multiscale phenomena**.



HODLR matrices

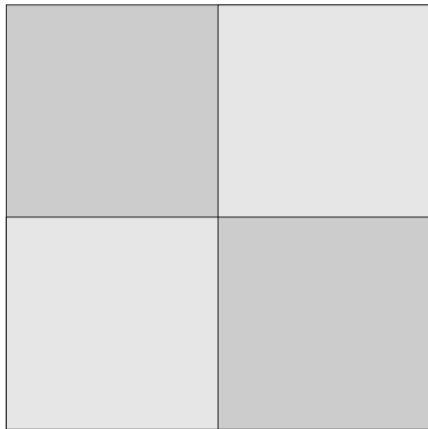


low-rank block



recursive block

HODLR matrices

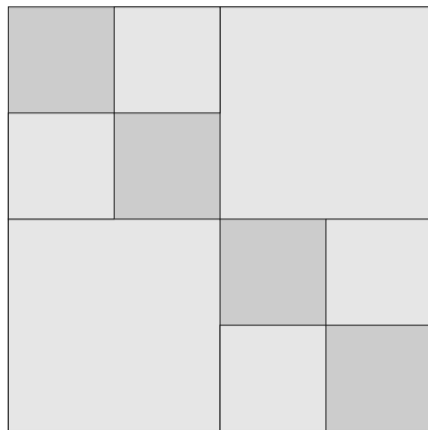


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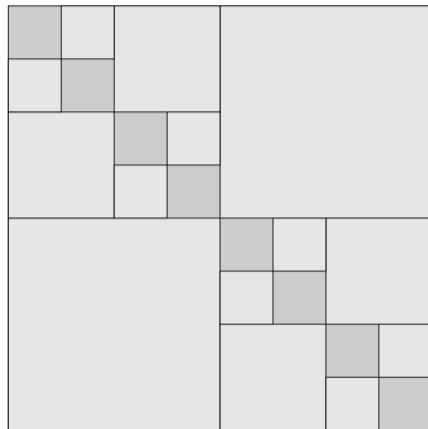


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HODLR Matrices

Definition. Fix a rank parameter k . We say a $n \times n$ matrix \mathbf{A} is HODLR(k) if $n \leq k$ or \mathbf{A} can be partitioned into $(n/2) \times (n/2)$ blocks

$$\mathbf{A} = \begin{bmatrix} \mathbf{A}_{1,1} & \mathbf{A}_{1,2} \\ \mathbf{A}_{2,1} & \mathbf{A}_{2,2} \end{bmatrix}$$

such that $\mathbf{A}_{1,2}$ and $\mathbf{A}_{2,1}$ are of rank at most k and $\mathbf{A}_{1,1}$ and $\mathbf{A}_{2,2}$ are each HODLR(k).

HODLR matrices have $O(kn \log(n))$ parameters.

There are several matvec algorithms for the recovery problem.³

³Lin, Lu, and Ying 2011; Martinsson 2016; Levitt and Martinsson 2022; Halikias and Townsend 2023.

Low-rank approximation from matrix-vector products

The Randomized SVD (RSVD) is a well-known algorithm for obtaining a low-rank approximation to a matrix \mathbf{B} :

1. Sample Gaussian matrix $\mathbf{\Omega}$
2. Form $\mathbf{Q} = \text{orth}(\mathbf{B}\mathbf{\Omega})$
3. Compute $\mathbf{X} = \mathbf{Q}^\top \mathbf{B}$ (minimize: $\|\mathbf{B} - \mathbf{Q}\mathbf{X}\|_F$)
4. Output $\mathbf{Q}[\mathbf{X}]_k$

Theorem. If \mathbf{B} is rank- k , and $\mathbf{\Omega}$ has $O(k)$ columns, then $\mathbf{Q}[\mathbf{X}]_k = \mathbf{B}$ (a.s.).

Peeling: an algorithm for the recovery problem⁴

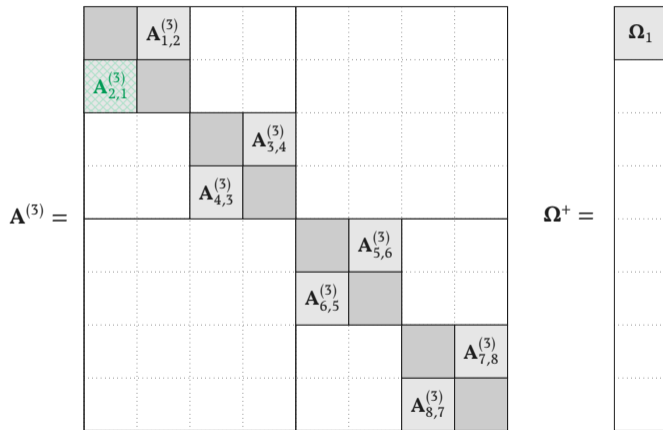
The algorithm works from the top layer down.

At each level, we **simultaneously** apply the RSVD to the low-rank off-diagonal blocks.

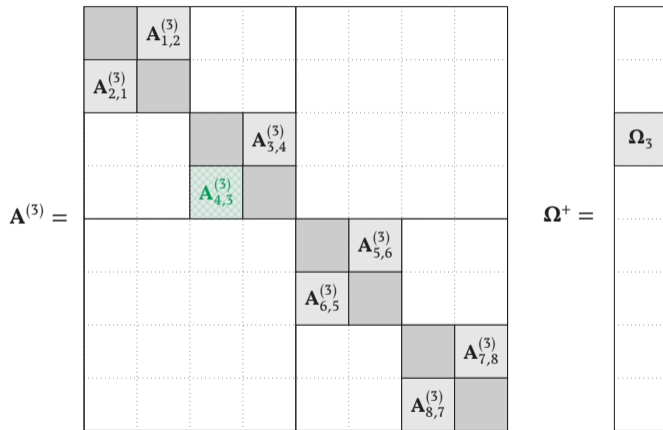
We then “peel” off these blocks before proceeding to the next level

⁴Lin, Lu, and Ying 2011; Martinsson 2016.

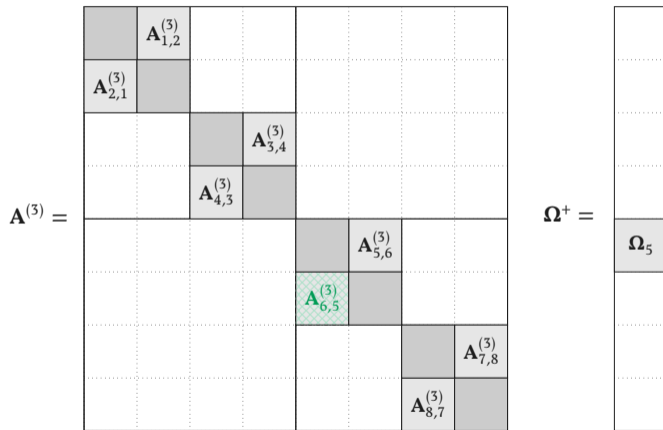
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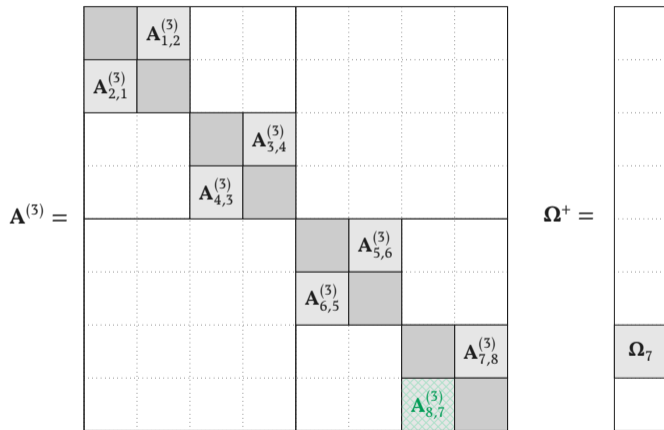
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$$\mathbf{A}^{(3)} = \begin{array}{|c|c|c|c|c|c|c|c|} \hline \text{ } & \mathbf{A}_{1,2}^{(3)} & \text{ } & \text{ } & \text{ } & \text{ } & \text{ } & \text{ } \\ \hline \mathbf{A}_{2,1}^{(3)} & \text{ } & \text{ } & \text{ } & \text{ } & \text{ } & \text{ } & \text{ } \\ \hline \text{ } & \text{ } & \text{ } & \mathbf{A}_{3,4}^{(3)} & \text{ } & \text{ } & \text{ } & \text{ } \\ \hline \text{ } & \text{ } & \mathbf{A}_{4,3}^{(3)} & \text{ } & \text{ } & \text{ } & \text{ } & \text{ } \\ \hline \text{ } & \text{ } & \text{ } & \text{ } & \text{ } & \mathbf{A}_{5,6}^{(3)} & \text{ } & \text{ } \\ \hline \text{ } & \text{ } & \text{ } & \text{ } & \mathbf{A}_{6,5}^{(3)} & \text{ } & \text{ } & \text{ } \\ \hline \text{ } & \text{ } & \text{ } & \text{ } & \text{ } & \text{ } & \mathbf{A}_{7,8}^{(3)} & \text{ } \\ \hline \text{ } & \text{ } & \text{ } & \text{ } & \text{ } & \mathbf{A}_{8,7}^{(3)} & \text{ } & \text{ } \\ \hline \end{array} \quad \mathbf{\Omega}^+ = \begin{array}{|c|} \hline \mathbf{\Omega}_1 \\ \hline \text{ } \\ \hline \mathbf{\Omega}_3 \\ \hline \text{ } \\ \hline \mathbf{\Omega}_5 \\ \hline \text{ } \\ \hline \mathbf{\Omega}_7 \\ \hline \text{ } \\ \hline \end{array}$$

From $\mathbf{A}^{(3)}\mathbf{\Omega}^+$ we get sketches: $\mathbf{A}_{2,1}^{(3)}\mathbf{\Omega}_1$, $\mathbf{A}_{4,3}^{(3)}\mathbf{\Omega}_3$, $\mathbf{A}_{6,5}^{(3)}\mathbf{\Omega}_5$, $\mathbf{A}_{8,7}^{(3)}\mathbf{\Omega}_7$.

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Peeling: an algorithm for the recovery problem

At each level we use $O(k)$ matrix-vector products with \mathbf{A} and \mathbf{A}^\top .

There are $\log_2(n/k) \leq \log_2(n)$ levels until the blocks are of size k

- then we can directly recover them at once with k products

Theorem. We can recover a HODLR matrix using $O(k \log_2(n))$ matvecs.

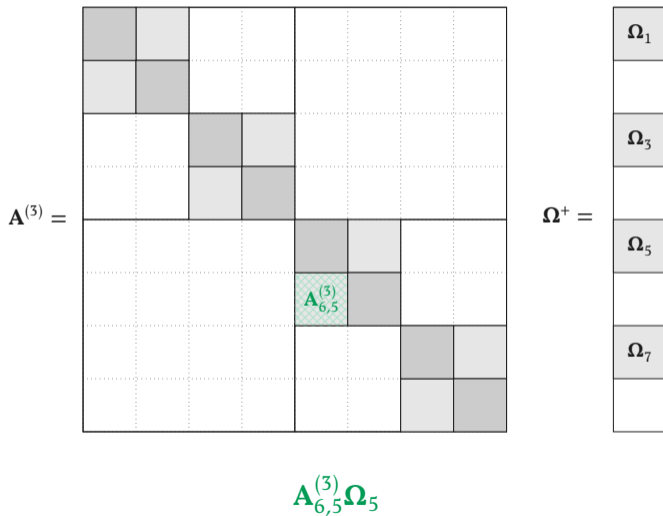
Peeling with error?

A variant of the peeling algorithm can be used to approximate the solution operator of elliptic PDEs (2024 SIAM Linear Algebra Best Paper Prize winner).⁵

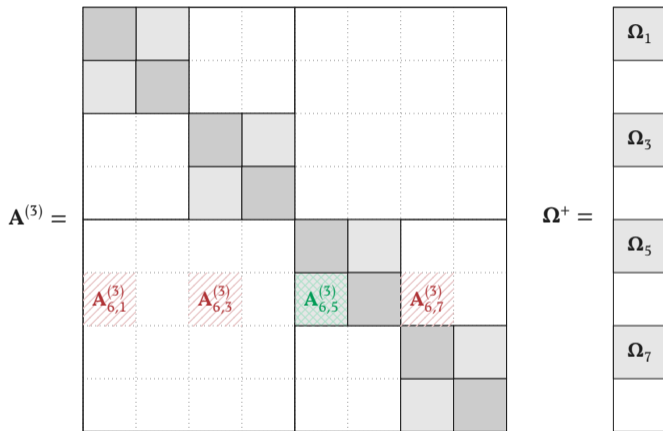
Boullé and Townsend 2022: Is there a peeling-type algorithm that works for nearly-HODLR matrices?

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Does peeling work on non-HODLR matrices?



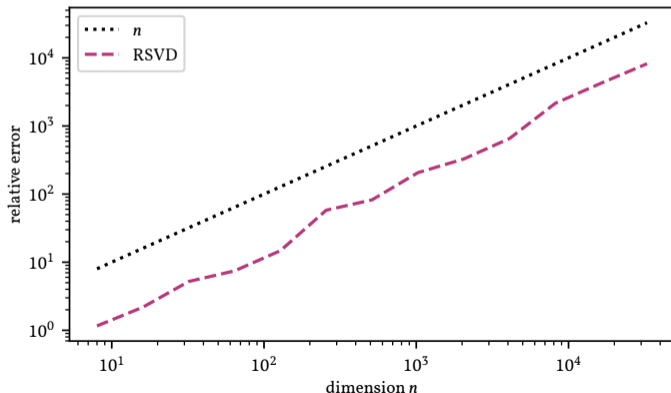
Does peeling work on non-HODLR matrices?



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Does peeling work on non-HODLR matrices?

If all the error at a level can propagate to the next level, then the total error doubles at each level. Exponential blow-up in the number of levels (linear in n)!



The HODLR approximation problem⁶

Problem. Given an $n \times n$ matrix \mathbf{A} , accessible only by matrix-vector products, a rank parameter k , and an accuracy parameter ε , find a HODLR(k) matrix $\tilde{\mathbf{A}}$ such that

$$\|\mathbf{A} - \tilde{\mathbf{A}}\|_{\text{F}} \leq (1 + \varepsilon) \min_{\mathbf{H} \in \text{HODLR}(k)} \|\mathbf{A} - \mathbf{H}\|_{\text{F}}.$$

Theorem. There is an efficient matvec algorithm for HODLR approximation.

Note: The best HODLR approximation to \mathbf{A} is obtained by applying a rank- k SVD to each low-rank block of \mathbf{A} .

- This is too expensive in the matrix-vector product model (n products)

⁶Chen et al. 2025.

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A perturbation bound for the RSVD

We prove a perturbation bound for the RSVD. This is likely of independent interest.

Theorem. Let $\mathbf{Q} = \text{orth}(\mathbf{B}\mathbf{\Omega} + \mathbf{E}_1)$ and $\mathbf{X} = \mathbf{Q}^\top \mathbf{B} + \mathbf{E}_2$. Then

$$\|\mathbf{B} - \mathbf{Q}[\mathbf{X}]_k\|_F \leq \underbrace{\|\mathbf{E}_1 \mathbf{\Omega}_{\text{top}}^\dagger\|_F + 2\|\mathbf{E}_2\|_F}_{\text{perturbations}} + \underbrace{\left(\|\mathbf{\Sigma}_{\text{bot}}\|_F^2 + \|\mathbf{\Sigma}_{\text{bot}} \mathbf{\Omega}_{\text{bot}} \mathbf{\Omega}_{\text{top}}^\dagger\|_F^2\right)^{1/2}}_{\text{classical RSVD bound}}.$$

When $\mathbf{\Omega}$ has $O(k/\varepsilon)$ columns, $\mathbf{\Omega}_{\text{top}}$ is a $k \times O(k/\varepsilon)$ Gaussian matrix which has a small pseudoinverse:

$$\mathbb{E}[(\mathbf{\Omega}_{\text{top}}^\dagger)^\top \mathbf{\Omega}_{\text{top}}^\dagger] = \mathbb{E}[(\mathbf{\Omega}_{\text{top}} \mathbf{\Omega}_{\text{top}}^\top)^{-1}] = \varepsilon \mathbf{I}.$$

Takeaway: The pseudoinverse will help damp the perturbation \mathbf{E}_1 , but (unsurprisingly) all of the perturbation \mathbf{E}_2 can propagate.

The RSVD tries to compute $\mathbf{Q}^\top \mathbf{B}$ directly; this is the solution to:

$$\min_{\mathbf{X}} \|\mathbf{A} - \mathbf{Q}\mathbf{X}\|_F.$$

Instead, we can solve a sketched problem:

$$\min_{\mathbf{X}} \|\Psi^\top \mathbf{A} - \Psi^\top \mathbf{Q}\mathbf{X}\|_F.$$

This means $\mathbf{X} = (\Psi^\top \mathbf{Q})^\dagger \Psi^\top \mathbf{A}$.

Observation. By adding columns to Ψ , we can damp errors in the product $\Psi^\top \mathbf{A}$.

The algorithm is also non-adaptive (we can do products with Ψ in advance)

⁷Clarkson and Woodruff 2009; Tropp, Yurtsever, Udell, and Cevher 2017; Nakatsukasa 2020.

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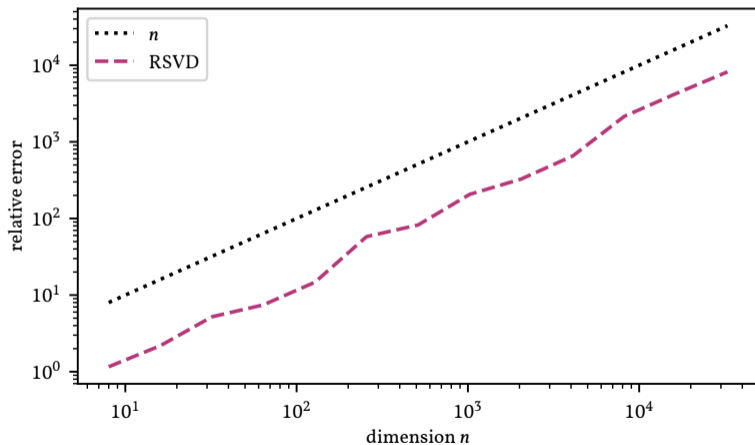
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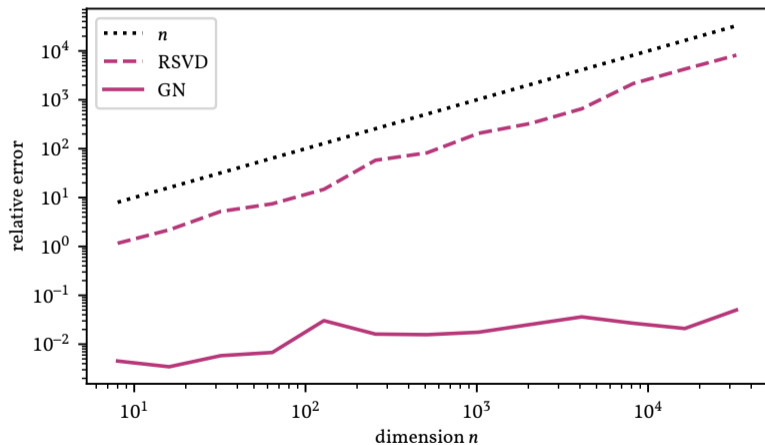
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Back to the hard instance



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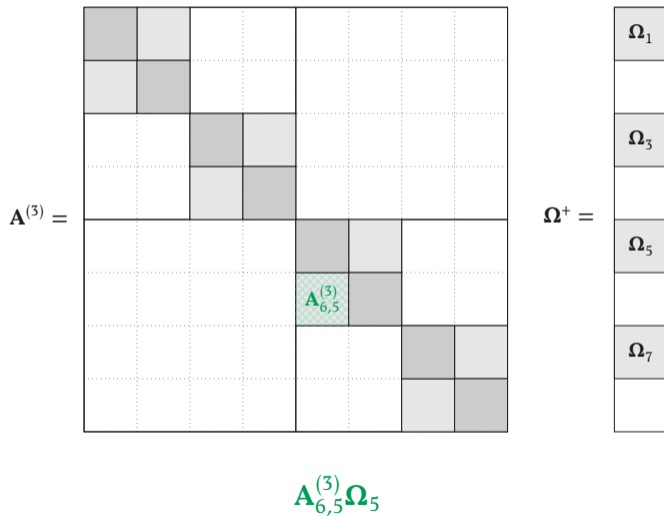


Another approach: perforated sketches

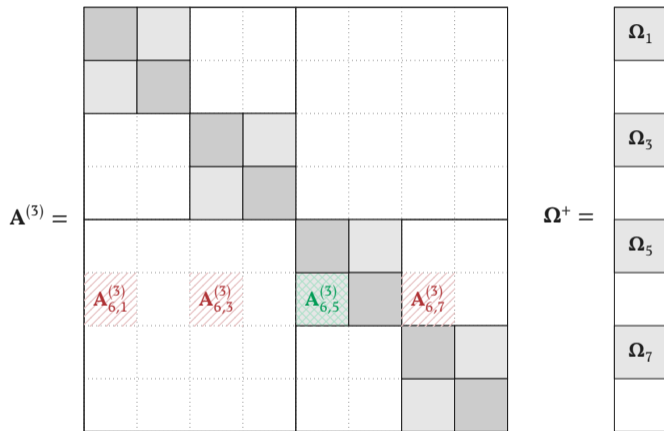
Because of the structure of peeling, the error happens when blocks of our sketch hit the error from our approximation of low-rank blocks at previous levels.

What if we just reduce how often this happens?

Perforated Block CountSketch

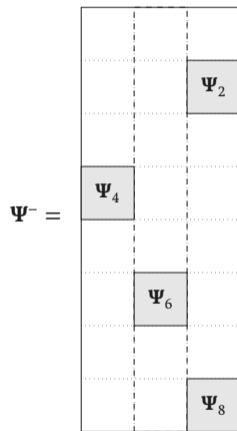
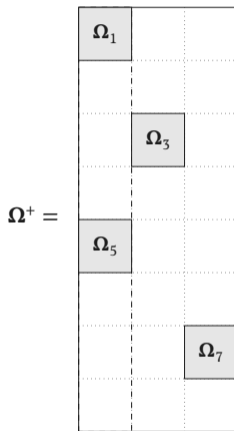
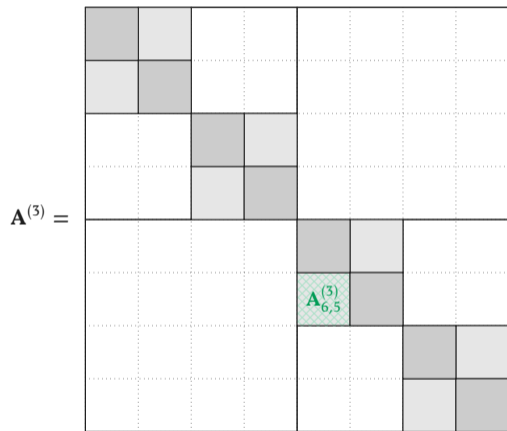


Perforated Block CountSketch



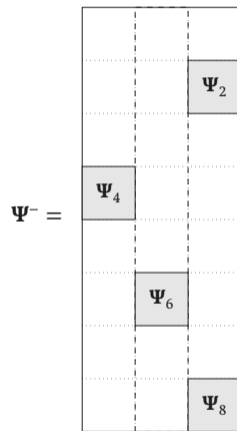
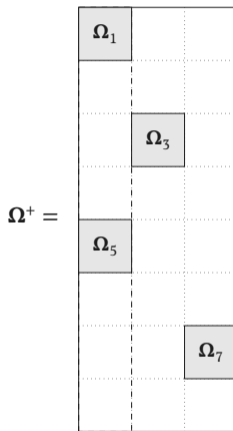
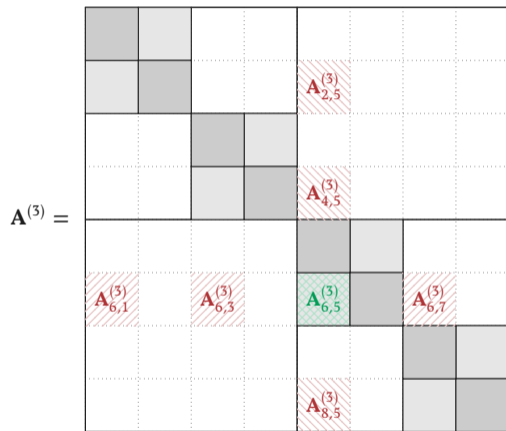
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Another idea: perforated Block CountSketch



$$\mathbf{A}_{6,5}^{(3)} \Omega_5$$

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$$\mathbf{A}_{6,5}^{(3)} \mathbf{\Omega}_5 + \mathbf{A}_{6,1}^{(3)} \mathbf{\Omega}_1$$

Our main result

Theorem. There exist matvec algorithms which use $O(k \log(n)/\beta^3)$ products with \mathbf{A} to obtain a $\text{HODLR}(k)$ matrix $\widetilde{\mathbf{A}}$ satisfying

$$\|\mathbf{A} - \widetilde{\mathbf{A}}\|_F \leq (1 + \beta)^{\log_2(n)} \min_{\mathbf{H} \in \text{HODLR}(k)} \|\mathbf{A} - \mathbf{H}\|_F.$$

Corollary. $(1 + \varepsilon)$ -optimal approximation with $O(k \log(n)^4/\varepsilon^3)$ matvecs

Corollary. $n^{0.01}$ -optimal approximation with $O(k \log(n))$ matvecs

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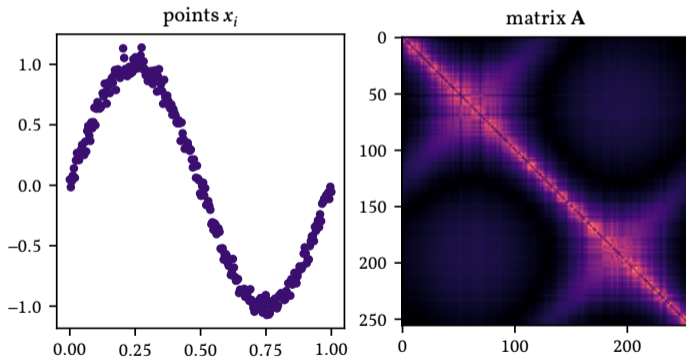
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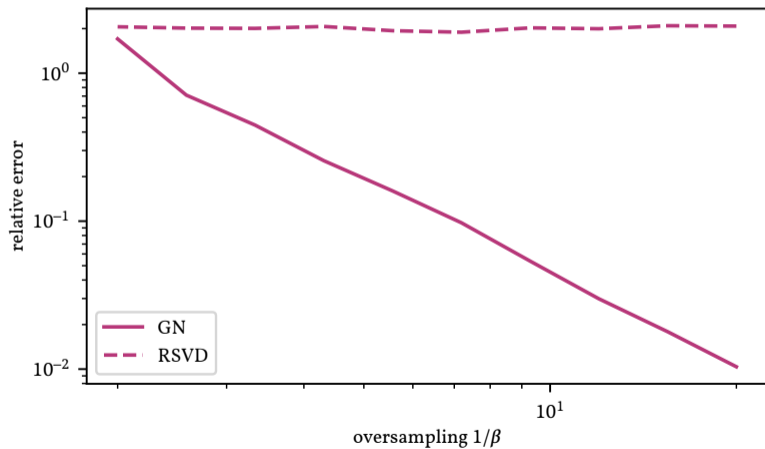
Corollary. $n^{0.01}$ -optimal approximation with $O(k \log(n))$ matvecs

Another experiment

Given points $x_i \in \mathbb{R}^2$, define $[\mathbf{A}]_{i,j} = -\log(\|x_i - x_j\|)$



Another experiment



Lower bounds?

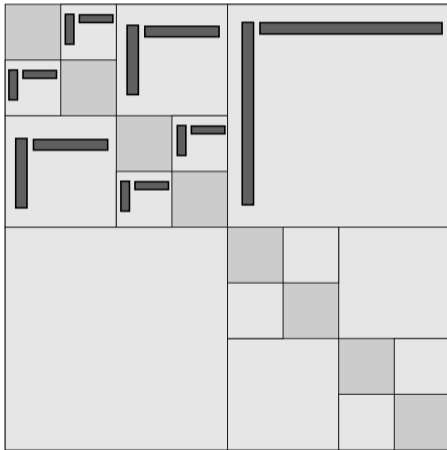
The matrix-vector query model often lets us prove **lower-bounds** against any matvec algorithm for a given task; i.e. study the **complexity** of a task.

This provides a very different approach for understanding how good algorithms are (compared to classical numerical analysis).

Theorem. There is a constant $C > 0$ such that for any k, n, ε , there exists a matrix A such that getting a $(1 + \varepsilon)$ -optimal HODLR approximation requires at least $C(k \log_2(n/k) + k/\varepsilon)$ matvecs.

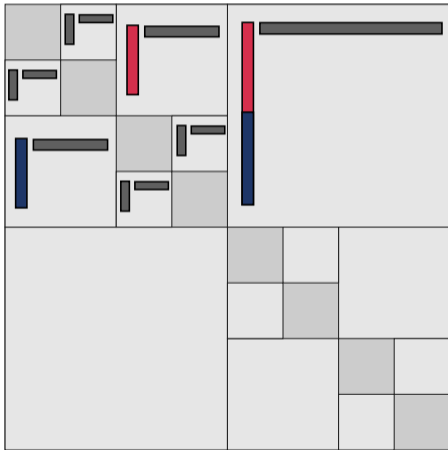
HSS matrices

The low-rank blocks of HSS matrices are related: $O(nk)$ parameters.



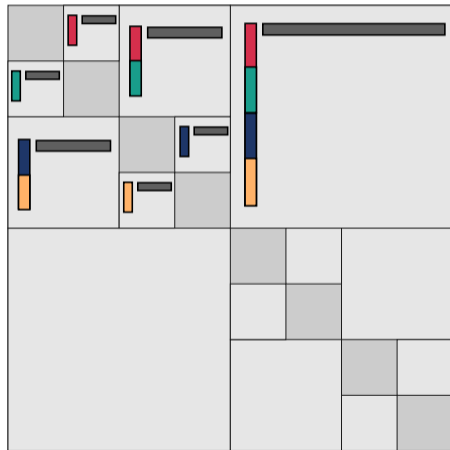
HSS matrices

The low-rank blocks of HSS matrices are related: $O(nk)$ parameters.



HSS matrices

The low-rank blocks of HSS matrices are related: $O(nk)$ parameters.



HSS is tricky!

Many papers study HSS recovery.⁸

The nestedness of column-spaces across levels adds lots of relations which make the approximation problem much harder.

- No known polynomial algorithm known for constant factor HSS approximation?!
- In fact, not even clear what to do in exponential time.

We prove:

Theorem. Can get $O(\log(n))$ -optimal HSS approximation in $O(kn^2)$ time.

⁸Xia, Chandrasekaran, Gu, and Li 2010; Levitt and Martinsson 2022; Halikias and Townsend 2023.

What's next?

Big goal: general theory for structured matrix approximation problem

- Correct $\log(n)$ and ε rates for the algorithms we study?
 - Limited by the best known bounds for Generalized Nyström: $O(k/\varepsilon^3)$
- True stability analysis (e.g. for floating point arithmetic)
 - Working on with students at NYU
- Adaptive algorithms
- Other hierarchical classes?
- Better understanding of (non-adaptive) low-rank approximation

Generalized Nyström (perturbation) analysis

Extend \mathbf{Q} to an orthogonal matrix $[\mathbf{Q} \ \widehat{\mathbf{Q}}]$, and write $\boldsymbol{\Psi}_1 = \boldsymbol{\Psi}^\top \mathbf{Q}$ and $\boldsymbol{\Psi}_2 = \boldsymbol{\Psi}^\top \widehat{\mathbf{Q}}$.

By orthogonal invariance, $\boldsymbol{\Psi}_1$ and $\boldsymbol{\Psi}_2$ are independent Gaussian matrices!

First observe:

$$\boldsymbol{\Psi}^\top \mathbf{B} = \boldsymbol{\Psi}^\top (\mathbf{Q}\mathbf{Q}^\top + \widehat{\mathbf{Q}}\widehat{\mathbf{Q}}^\top) \mathbf{B} = \boldsymbol{\Psi}_1 \mathbf{Q}^\top \mathbf{B} + \boldsymbol{\Psi}_2 \widehat{\mathbf{Q}}^\top \mathbf{B}.$$

Therefore:

$$\mathbf{X} = (\boldsymbol{\Psi}^\top \mathbf{Q})^\dagger (\boldsymbol{\Psi}^\top \mathbf{B}) = \boldsymbol{\Psi}_1^\dagger \boldsymbol{\Psi}_1 \mathbf{Q}^\top \mathbf{B} + \boldsymbol{\Psi}_1^\dagger \boldsymbol{\Psi}_2 \widehat{\mathbf{Q}}^\top \mathbf{B} = \mathbf{Q}^\top \mathbf{B} + \boldsymbol{\Psi}_1^\dagger \boldsymbol{\Psi}_2 \widehat{\mathbf{Q}}^\top \mathbf{B}.$$

Adding more columns to $\boldsymbol{\Psi}$ (and hence $\boldsymbol{\Psi}_1$) reduces the error in the second term.

Generalized Nyström (perturbation) analysis

Extend \mathbf{Q} to an orthogonal matrix $[\mathbf{Q} \ \widehat{\mathbf{Q}}]$, and write $\boldsymbol{\Psi}_1 = \boldsymbol{\Psi}^\top \mathbf{Q}$ and $\boldsymbol{\Psi}_2 = \boldsymbol{\Psi}^\top \widehat{\mathbf{Q}}$.

By orthogonal invariance, $\boldsymbol{\Psi}_1$ and $\boldsymbol{\Psi}_2$ are independent Gaussian matrices!

First observe:

$$\boldsymbol{\Psi}^\top \mathbf{B} + \mathbf{E} = \boldsymbol{\Psi}^\top (\mathbf{Q}\mathbf{Q}^\top + \widehat{\mathbf{Q}}\widehat{\mathbf{Q}}^\top) \mathbf{B} + \mathbf{E} = \boldsymbol{\Psi}_1 \mathbf{Q}^\top \mathbf{B} + \boldsymbol{\Psi}_2 \widehat{\mathbf{Q}}^\top \mathbf{B} + \mathbf{E}.$$

Therefore:

$$\mathbf{X} = (\boldsymbol{\Psi}^\top \mathbf{Q})^\dagger (\boldsymbol{\Psi}^\top \mathbf{B} + \mathbf{E}) = \boldsymbol{\Psi}_1^\dagger \boldsymbol{\Psi}_1 \mathbf{Q}^\top \mathbf{B} + \boldsymbol{\Psi}_1^\dagger \boldsymbol{\Psi}_2 \widehat{\mathbf{Q}}^\top \mathbf{B} + \boldsymbol{\Psi}_1^\dagger \mathbf{E} = \mathbf{Q}^\top \mathbf{B} + \boldsymbol{\Psi}_1^\dagger \boldsymbol{\Psi}_2 \widehat{\mathbf{Q}}^\top \mathbf{B} + \boldsymbol{\Psi}_1^\dagger \mathbf{E}.$$

Adding more columns to $\boldsymbol{\Psi}$ (and hence $\boldsymbol{\Psi}_1$) reduces the error in the second term.

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